Partial Fractions Expansion of Rational Functions

An Application of the Fundamental Theorem of Algebra William J. Martin, WPI

The Fundamental Theorem of Algebra is important throughout mathematics. Here are two equivalent statements of this theorem.

Theorem 1 (Fundamental Theorem of Algebra, Building Block). Every non-constant polynomial with complex coefficients has a complex root.

Theorem 2 (Fundamental Theorem of Algebra, Factorization). A non-zero polynomial of degree n with complex coefficients has exactly n complex roots, counting multiplicities.

Let $\mathbb{C}[z]$ denote the ring of polynomials with complex coefficients. Each element $f(z) \in \mathbb{C}[z]$ has the form

$$f(z) = a_n z^n + a_{n-1} z^{n-1} + \dots + a_1 z + a_0$$

for some complex numbers a_0, a_1, \ldots, a_n . Provided $a_n \neq 0$, we say f(z) has degree n. We say f(z) is a *monic* polynomial if $a_n = 1$ (the leading coefficient is equal to one).

The first version of the FTA states that, if $f(z) \in \mathbb{C}[z]$ has degree $n \geq 1$, then there is some complex number $c \in \mathbb{C}$ such that

$$f(c) = a_n c^n + a_{n-1} c^{n-1} + \dots + a_1 c + a_0 = 0.$$

By the Factor Theorem, this tells us that the polynomial z - c divides f(z). So there is some monic polynomial g(z) such that

$$f(z) = a_n(x - c)g(z).$$

Since \mathbb{C} is an integral domain (in fact, it's a field), we know that g(z) has degree n-1. Applying the Principle of Mathematical Induction, we may assume that g(z) has exactly n-1 complex roots. So the "building block" version of the FTA gives us the factorization version of the FTA:

For any polynomial $f(z) = a_n z^n + a_{n-1} z^{n-1} + \cdots + a_1 z + a_0$ with complex coefficients, with $n \ge 0$ and $a_n \ne 0$, there exist n complex numbers

$$r_1, r_2, \ldots, r_n$$

not necessarily distinct, such that f(z) factors into irreducibles as

$$f(z) = a_n(z - r_1)(z - r_2) \cdots (z - r_n).$$

The Fundamental Theorem of Algebra quickly implies (and it is essentially equivalent to) the following useful theorem about polynomials with real coefficients. The ring $\mathbb{R}[x]$ consists of all polynomials with real coefficients, that is, all

$$f(x) = a_n x^n + a_{n-1} x^{n-1} + \dots + a_1 x + a_0$$

where $n \ge 0$ and all a_i are real numbers.

Quadratic polynomials $f(x) = ax^2 + bx + c$ fall into two classes: those with real roots¹ (which happens when $b^2 - 4ac \ge 0$) and those with no real roots (which happens when $b^2 - 4ac < 0$). These latter type are called "irreducible quadratics".

Theorem 3 (Fundamental Theorem of Algebra, Real Coefficients). Every nonzero polynomial with real coefficients factors uniquely into linear polynomials and irreducible quadratics.

Let's unwrap this and spell out exactly what it tells us.

If $f(x) = a_n x^n + a_{n-1} x^{n-1} + \dots + a_1 x + a_0$ with $n \ge 0$ and $a_n \ne 0$, then there are unique (up to ordering) monic polynomials $\ell_1(x) = x - r_1$, $\ell_2(x) = x - r_2$, $\dots \ell_s(x) = x - r_s$, and $q_1(x) = x^2 + b_1 x + c_1$, $q_2(x) = x^2 + b_2 x + c_2$, $\dots q_t(x) = x^2 + b_t x + c_t$ with each $b_j^2 - 4c_j < 0$ (and, comparing degrees, s + 2t = n) such that

$$f(x) = a_n(x - r_1)(x - r_2) \cdots (x - r_s)(x^2 + b_1x + c_1)(x^2 + b_2x + c_2) \cdots (x^2 + b_tx + c_t).$$

We then see that this polynomial has exactly s real roots (namely, r_1, r_2, \ldots, r_s) and 2t complex roots that are not real:

$$\frac{-b_1}{2} + \frac{i}{2}\sqrt{4c_1 - b_1^2}, \ \frac{-b_1}{2} - \frac{i}{2}\sqrt{4c_1 - b_1^2}, \ \dots, \frac{-b_t}{2} + \frac{i}{2}\sqrt{4c_t - b_t^2}, \ \frac{-b_t}{2} - \frac{i}{2}\sqrt{4c_t - b_t^2}.$$

To prove this result using the FTA, we observe that complex conjugation (sending z = a + bi to $\overline{z} = a - bi$) is a ring isomorphism (one-to-one and onto) from the field of complex numbers to itself:

$$\overline{z+w} = \overline{z} + \overline{w}, \qquad \overline{z \cdot w} = \overline{z} \cdot \overline{w},$$

with $\overline{z} = z$ if and only if z is a real number.

Suppose that $f(x) = a_n x^n + a_{n-1} x^{n-1} + \cdots + a_1 x + a_0$ is a polynomial with real coefficients and suppose that the complex number z is a root of f. Then

$$0 = f(z) = a_n z^n + a_{n-1} z^{n-1} + \dots + a_1 z + a_0$$

and, applying complex conjugation to both sides,

$$0 = \bar{0} = \overline{a_n z^n + a_{n-1} z^{n-1} + \dots + a_1 z + a_0} = a_n \bar{z}^n + a_{n-1} \bar{z}^{n-1} + \dots + a_1 \bar{z} + a_0$$

showing that \bar{z} is also a root of f(x). So the non-real roots of f(x) come in conjugate pairs.

It is important to remember that, for any complex number z, both $z + \bar{z}$ and $z \cdot \bar{z}$ are real numbers. (Please check this.) So the polynomial $(x - z)(x - \bar{z})$ has the form $x^2 + bx + c$ where b and c are real numbers.

What does this achieve for us? We are assuming that f(x) has real coefficients. We apply the Fundamental Theorem of Algebra to express $f(x) = a_n x^n + \cdots + a_1 x + a_0$ as

$$f(x) = a_n(x - r_1)(x - r_2) \cdots (x - r_n)$$

¹Note that the quadratic $f(x) = ax^2 + bx + c$ has a double root (or repeated root) precisely when $b^2 - 4ac = 0$.

where r_1, \ldots, r_n are the *n* complex roots of *f*. We can re-order these so that the real roots come first: say r_1, \ldots, r_s are real numbers and all the rest are non-real complex numbers. We just learned that these come in conjugate pairs. So n - s is even and we can re-order terms so that r_{s+2} is the conjugate of r_{s+1} and r_{s+4} is the conjugate of r_{s+3} and so on. But then each pair of linear factors $(x - r_{s+j-1})(x - r_{s+j})$ where *j* is even expands into an irreducible quadratic polynomial with real coefficients: since $r_{s+j-1} = \bar{r}_{s+j}$, we have

$$(x - r_{s+j-1})(x - r_{s+j}) = x^2 - (r_{s+j} + \bar{r}_{s+j})x + r_{s+j}\bar{r}_{s+j}.$$

This is the origin of the irreducible quadratic factors in the "real" version of the FTA.

1 Partial Fractions Expansion

The ring of rational functions contains all ratios of two polynomials in x

$$\left\{\frac{f(x)}{g(x)} \mid f(x), g(x) \in \mathbb{R}[x], \ g(x) \text{ not the zero polynomial} \right\}$$

is a commutative ring with operations

$$\frac{f(x)}{g(x)} + \frac{h(x)}{k(x)} = \frac{f(x)k(x) + h(x)g(x)}{g(x)k(x)} \quad \text{and} \quad \frac{f(x)}{g(x)} \cdot \frac{h(x)}{k(x)} = \frac{f(x)h(x)}{g(x)k(x)}$$

But this happens to be a field! Every non-zero element has an inverse. Let's check. Suppose r is a rational function which is not the zero element. Then r takes the form $\frac{f(x)}{g(x)}$ for some polynomials f(x) and g(x) and r being non-zero means that f(x) is not the zero polynomial. So the ratio $\frac{g(x)}{f(x)}$ is also a rational function and it is easy to check that this is the inverse of $\frac{f(x)}{f(x)}$.

 $\frac{f(x)}{g(x)}$. Since 1 is a non-zero polynomial, we see that every polynomial is a rational function: write f(x) as $\frac{f(x)}{1}$. (In fact, the identity element 1 for this field is technically written $\frac{1}{1}$.) By long division (the Division Algorithm for polynomials), we can express any rational function f(x)/g(x) in the form q(x) + r(x)/g(x) where the degree of r is less than the degree of g. For there is a unique quotient q(x) and a unique remainder r(x) such that

- f(x) = g(x)q(x) + r(x)
- either r(x) is the zero polynomial or deg $r < \deg g$

What we are after is a way to reverse the common denominator approach to adding fractions. We can all easily compute

$$\frac{3}{7} + \frac{2}{5} = \frac{29}{35}$$

but how do we find small integers a and b such that

$$\frac{17}{33} = \frac{a}{3} + \frac{b}{11}$$
?

Do we even know that such a and b exist? We can multiply both sides by 33 to get

$$17 = 11a + 3b.$$

Since 11 and 3 are relatively prime, we can find integers x and y such that

$$11x + 3y = 1.$$

For example x = -1, y = 4 works. We can then multiply both sides of this equation by 17 to find

$$17 = 11(-17) + 3(68)$$

but the expansion

$$\frac{17}{33} = -\frac{17}{3} + \frac{68}{11}$$

doesn't seem like much of a simplification! We'd rather have $\frac{17}{33} = \frac{1}{3} + \frac{2}{11}$. [I'll leave this mystery to you to resolve.]

We want to do the same thing for polynomials: "undo" the common denominator calculation. For instance, while we all can compute

$$\frac{3}{x+5} - \frac{2}{x-5} = \frac{x-25}{x^2-25}$$

we must solve a system of linear equations to discover that

$$\frac{12}{x^2 - x - 2} = \frac{4}{x - 2} - \frac{4}{x + 1}$$

Example: Find real numbers A and B such that

$$\frac{4x+1}{x^2-3x-10} = \frac{A}{x+2} + \frac{B}{x-5}.$$

Solution: Multiply both sides by $x^2 - 3x - 10 = (x + 2)(x - 5)$ to find

$$4x + 1 = (A + B)x + (2B - 5A).$$

In order for these polynomials to be equal, all corresponding coefficients must be equal. So we obtain the linear system

$$A + B = 4, \qquad 2B - 5A = 1$$

which we quickly solve to find A = 1, B = 3. \Box

The fact that this can always be achieved follows from the following lemmas. Comparing to our handling of mixed numbers, the first lemma is analogous to saying that, if m and n are integers with $0 \le n < 7$ which satisfy

$$5\frac{2}{7} = m + \frac{n}{7},$$

then m = 5 and n = 2.

Lemma 1. If $f_1, g_1, h_1, f_2, g_2, h_2$ are all polynomials with deg $g_1 < \deg h_1$ and deg $g_2 < \deg h_2$ and the equation

$$f_1(x) + \frac{g_1(x)}{h_1(x)} = f_2(x) + \frac{g_2(x)}{h_2(x)}$$

holds, then $f_1(x) = f_2(x)$ and $h_2(x)g_1(x) = h_1(x)g_2(x)$.

<u>Proof:</u> Omitted. \Box

Lemma 2. If f(x) and g(x) are polynomials with greatest common divisor d(x), then there exist polynomials s(x) and t(x) satisfying the equation

$$f(x) \cdot s(x) + g(x) \cdot t(x) = d(x).$$

In particular, polynomials f and g have no common factor (i.e., the two polynomials are relatively prime) if and only if there exist polynomials s(x) and t(x) satisfying the equation $f(x) \cdot s(x) + g(x) \cdot t(x) = 1$.

<u>Proof:</u> This follows from the Extended Euclidean Algorithm much in the same way as the analogous theorem for integers. \Box

Lemma 3. Suppose we are given a rational function f(x)/g(x) where deg $f < \deg g$. If g(x) factors as $g(x) = p(x) \cdot q(x)$ where polynomials p(x) and q(x) are relatively prime, then there exist polynomials s(x) and t(x) such that

- $\frac{f(x)}{g(x)} = \frac{s(x)}{p(x)} + \frac{t(x)}{q(x)}$
- $\deg s < \deg p$ and $\deg t < \deg q$.

<u>Proof:</u> Since p(x) and q(x) have gcd equal to 1, there are polynomials S(x) and T(x) such that S(x)q(x) + T(x)p(x) = 1. Multiplying both sides by the rational function f(x)/g(x), we get

$$\frac{S(x)f(x)}{p(x)} + \frac{T(x)f(x)}{q(x)} = \frac{f(x)}{g(x)}$$
(1)

since $g(x) = p(x) \cdot q(x)$. Now it may be that one of the two rational functions on the left has a numerator with higher degree than its denominator. If deg $S + \deg f \ge \deg p$, then write

$$S(x)f(x) = p(x)u(x) + r(x)$$

where deg $r < \deg p$; otherwise, take u(x) = 0 and r(x) = S(x)f(x). Likewise, write

$$T(x)f(x) = q(x)w(x) + v(x)$$

for some unique quotient w(x) (possibly zero) and some remainder v(x) with deg $v < \deg q$. Then Equation (1) can be written

$$u(x) + \frac{r(x)}{p(x)} + w(x) + \frac{v(x)}{q(x)} = \frac{f(x)}{g(x)}$$

and by Lemma 1, we have u(x)+w(x) is the zero polynomial since deg $f < \deg g$. Eliminating these two terms gives us the desired decomposition. \Box

(You might now wish to go back to the equation $\frac{17}{33} = -\frac{17}{3} + \frac{68}{11}$ and see how this idea allows us to simplify it to $\frac{17}{33} = \frac{1}{3} + \frac{2}{11}$.

Lemma 4. Consider rational functions where the denominator is expressible as some irreducible polynomial in $\mathbb{R}[x]$ to some power.

(a) If deg f(x) < k and $c \in \mathbb{R}$, then there exist unique real numbers $a_k, a_{k-1}, \ldots, a_1$ such that

$$\frac{f(x)}{(x-c)^k} = \frac{a_k}{(x-c)^k} + \frac{a_{k-1}}{(x-c)^{k-1}} + \dots + \frac{a_2}{(x-c)^2} + \frac{a_k}{x-c}.$$

(b) If deg f(x) < 2k and $b, c \in \mathbb{R}$ with $b^2 - 4c < 0$, then there exist unique real numbers $A_k, A_{k-1}, \ldots, A_1$ and $B_k, B_{k-1}, \ldots, B_1$ such that

$$\frac{f(x)}{(x^2+bx+c)^k} = \frac{A_k x + B_k}{(x^2+bx+c)^k} + \frac{A_{k-1}x + B_{k-1}}{(x^2+bx+c)^{k-1}} + \dots + \frac{A_1x + B_1}{x^2+bx+c}$$

<u>Proof</u>: Repeatedly apply the Division Algorithm (long division) for polynomials: in case (a),

$$f(x) = (x - c)f_k(x) + a_k$$

so that

$$\frac{f(x)}{(x-c)^k} = \frac{a_k}{(x-c)^k} + \frac{f_k(x)}{(x-c)^{k-1}}$$

and write $f_k(x) = (x-c)f_{k-1}(x) + a_{k-1}$ so that $\frac{f_k(x)}{(x-c)^{k-1}} = \frac{a_{k-1}}{(x-c)^{k-1}} + \frac{f_{k-1}(x)}{(x-c)^{k-2}}$ and so on. \Box

Theorem 4. Let f(x)/g(x) be any rational function which is a quotient of polynomials with real coefficients. Factor g(x) as

$$g(x) = b_n (x - r_1)^{k_1} (x - r_2)^{k_2} \cdots (x - r_s)^{k_s} (x^2 + b_1 x + c_1)^{\ell_1} (x^2 + b_2 x + c_2)^{\ell_2} \cdots (x^2 + b_t x + c_t)^{\ell_t}.$$

Then there exists a unique polynomial H(x) and unique real numbers

$$A_{k_{1},1}, A_{k_{1}-1,1}, \dots, A_{1,1}, A_{k_{2},2}, \dots, A_{1,2}, \dots, A_{k_{s},s}, \dots, A_{1,s}, B_{\ell_{1},1}, B_{\ell_{1}-1,1}, \dots, B_{1,1}, B_{\ell_{2},2}, \dots, B_{1,2}, \dots, B_{\ell_{t},t}, \dots, B_{1,t}, C_{\ell_{1},1}, C_{\ell_{1}-1,1}, \dots, C_{1,1}, C_{\ell_{2},2}, \dots, C_{1,2}, \dots, C_{\ell_{t},t}, \dots, C_{1,t},$$

such that

$$\frac{f(x)}{g(x)} = H(x) + \frac{A_{k_{1},1}}{(x-r_{1})^{k_{1}}} + \frac{A_{k_{1}-1,1}}{(x-r_{1})^{k_{1}-1}} + \dots + \frac{A_{1,1}}{x-r_{1}} + \frac{A_{k_{2},2}}{(x-r_{2})^{k_{2}}} + \dots + \frac{A_{1,2}}{x-r_{2}} + \dots + \frac{A_{1,2}}{x-r_{2}} + \dots + \frac{A_{k_{s},s}}{(x-r_{s})^{k_{s}}} + \dots + \frac{A_{1,s}}{x-r_{s}} + \frac{B_{\ell_{1},1}x + C_{\ell_{1},1}}{(x^{2} + b_{1}x + c_{1})^{\ell_{1}}} + \frac{B_{\ell_{1}-1,1}x + C_{\ell_{1}-1,1}}{(x^{2} + b_{1}x + c_{1})^{\ell_{1}-1}} + \dots + \frac{B_{1,1}x + C_{1,1}}{x^{2} + b_{1}x + c_{1}} + \frac{B_{\ell_{2},2}x + C_{\ell_{2},2}}{(x^{2} + b_{2}x + c_{2})^{\ell_{2}}} + \dots + \frac{B_{1,2}x + C_{1,2}}{x^{2} + b_{2}x + c_{2}} + \dots + \frac{B_{\ell_{t},t}x + C_{\ell_{t},t}}{(x^{2} + b_{t}x + c_{t})^{\ell_{t}}} + \dots + \frac{B_{1,t}x + C_{1,t}}{(x^{2} + b_{t}x + c_{t})}$$

<u>Proof:</u> This horrendous expression is obtained by applying the previous lemmas. First, if the degree of f(x) is larger or equal to the degree of g(x), then we apply the Division Algorithm to get the quotient H(x) and a remainder of degree less than deg g. So, for the remainder of this proof, we assume the numerator has smaller degree than the denominator.

Whenever g(x) factors, we can either break it up into relatively prime polynomials and apply Lemma 3 to write the fraction f(x)/g(x) as a sum of two fractions (rational functions) whose denominators have smaller degrees OR we have only repeated factors: $g(x) = p(x)^k$ for some polynomial p(x) and some integer $k \ge 1$. We only get stuck here when p(x) is irreducible in the ring $\mathbb{R}[x]$. By the FTA, we know that p(x) is either a linear polynomial x - r or an irreducible quadratic polynomial $x^2 + bx + c$ with $b^2 - 4c < 0$. In this case, we apply Lemma 4 to decompose $f(x)/(p(x))^k$ into a sum of terms $A/(x-r)^j$ (if p(x) = x - r) or $(Bx+C)/(x^2+bx+c)^j$ (if p(x) is quadratic. This eventually produces the partial fractions expansion promised by the theorem. \Box

Of course, the general expression is horrendous, so we should see examples to understand all the notation.

Suppose $f(x) = x^6 - 3x^4 + 8x^3 + x$ and $g(x) = x(x-5)^3(x^2+x+1)^2$. Then there exist unique real numbers

$$A_{1,1}, A_{3,2}, A_{2,2}, A_{1,2}, B_{2,1}, C_{2,1}, B_{1,1}, C_{1,1}$$

such that

$$\frac{x^6 - 3x^4 + 8x^3 + x}{x(x-5)^3(x^2 + x + 1)^2} = \frac{A_{1,1}}{x} + \frac{A_{3,2}}{(x-5)^3} + \frac{A_{2,2}}{(x-5)^2} + \frac{A_{1,2}}{x-5} + \frac{B_{2,1}x + C_{2,1}}{(x^2 + x + 1)^2} + \frac{B_{1,1}x + C_{1,1}}{x^2 + x + 1} + \frac{A_{2,2}}{(x-5)^2} + \frac{A_{2,2}}$$

In order to find these numbers, we multiply through both sides by g(x) to remove all denominators and then collect like terms. Matching coefficients of the polynomial on the left to coefficients of the polynomial on the right, term by term, we obtain a system of linear equations which we then solve for the $A_{i,j}$, $B_{i,j}$ and $C_{i,j}$.

The ten exercises for Homework 5 are on the next page.

Exercises

- 1. RSA encryption is based on algebra in the ring \mathbb{Z}_n where n = pq is a product of two large distinct primes p and q. Security is compromised if p and/or q is leaked. So our hope for security rests on the assumption that factoring the integer n is very hard. The proof that decryption works depends on the integer $\phi(n)$ where ϕ is Euler's totient function.
 - (a) Show how to compute $\phi(n)$ using n, p and q.
 - (b) Using the quadratic formula, show how to compute p and q if given n and $\phi(n)$.
- 2. Factor $f(x) = x^4 + 81$ into irreducibles over \mathbb{R} .
- 3. If z is the complex number given by $z = Re^{i\theta}$ where R > 0 and θ are real numbers, what is the equivalent expression for the complex number z^n where n is a positive integer? What about \sqrt{z} ?
- 4. Write $\frac{1}{x-2} \frac{3}{x-3} + \frac{2}{x-4}$ as a rational function f(x)/g(x).
- 5. Compute the partial fractions expansion of the rational function $(x^2-7)/(x^2-7x+10)$ and note that, in this case the polynomial H(x) is not zero.
- 6. Compute the partial fractions expansion of the rational function $(x^2 + 2x)/(x^3 1)$.
- 7. Compute the partial fractions expansion of $(x^4 x + 1)/[x^2(x^2 + 1)^2]$. Be careful!
- 8. Derive the quadratic formula by solving each of the following quadratic equations for x.
 - (a) $ax^2 = n$ (b) $x^2 + 2kx + k^2 = n$
 - (c) $a(x+k)^2 = n$
 - (d) $(x + \frac{b}{2a})^2 = \frac{b^2 4ac}{4a^2}$
- 9. Write down the first five rows of Pascal's Triangle and, beside this, write down the expansion of the polynomials $(x + y)^1$, $(x + y)^2$, $(x + y)^3$ and $(x + y)^4$.
- 10. Let $f(x) = x^n$ and let a be a real number. Show that

$$f(x+a) - f(a) = x \cdot \left[nx^{n-1} + \binom{n}{2} x^{n-2} + \dots + \binom{n}{n-1} \right].$$

You may use the Binomial Theorem in your derivation without proving it.