



- Ordinary

- Replace

- Double Down

1. (4 pts) Given a grammar G :

$$\begin{aligned}
 G : S &\rightarrow AB \mid C \\
 A &\rightarrow AC \mid AB \mid BC \mid BD \\
 B &\rightarrow BC \mid BD \mid a \mid bb \\
 C &\rightarrow CC \mid aa \mid bbb \\
 D &\rightarrow aa \mid bb
 \end{aligned}$$

Express as an equivalent grammar with no left recursion.

♣ A , B and C have left recursion, so we introduce R_A , R_B , and R_C .

$$\begin{aligned}
 G' : S &\rightarrow AB \mid C \\
 A &\rightarrow BC \mid BD \mid BCR_A \mid BDR_A \\
 R_A &\rightarrow C \mid B \mid CR_A \mid BR_A \\
 B &\rightarrow a \mid bb \mid aR_B \mid bbR_B \\
 R_B &\rightarrow C \mid D \mid CR_B \mid DR_B \\
 C &\rightarrow aa \mid bbb \mid aaR_C \mid bbbR_C \\
 R_C &\rightarrow C \mid CR_C \\
 D &\rightarrow aa \mid bb
 \end{aligned}$$

2. (4 pts) Give a regular expression for the language $L \subseteq \{a, b, c\}^*$ consisting of all strings which contain the substring $abab$ as well as the substring $baba$.

♣ There are many ways to do this. The two subwords to be included can occur in either order and can intersect in either 1 or 3 characters, in which case the string contains $ababa$ or $babab$ which is sufficient to insure the string is in the language. So

$$(a \cup b \cup c)^* [ababa \cup babab \cup abab(a \cup b \cup c)^* baba \cup baba(a \cup b \cup c)^* abab] (a \cup b \cup c)^*$$

3. (2 pts) On the back of this page, draw the state diagram of an automaton, deterministic, or non-deterministic, which expresses the language in problem 2.

♣ Is much easier to use a non-deterministic machine following the idea of problem 2.

