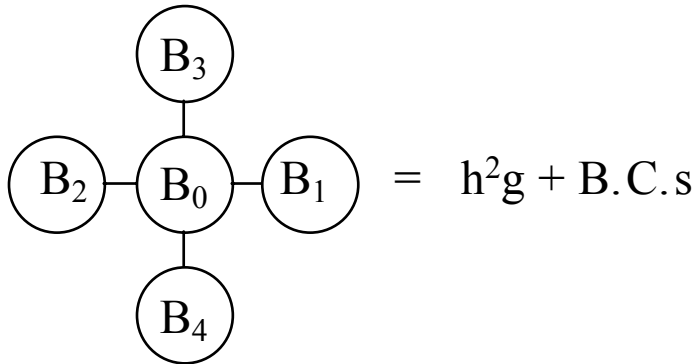


Last Time - Point Iterative Methods

$$[A] \{U\} = \{v\} \quad \text{for the system} \quad \nabla^2 U = g$$

with a computational molecule



[A] partitioned into [R] + [D] + [S]
 Below Diagonal Above

$$U_{i,j}^{n+1} = \frac{\omega}{B_0} \left[B_1 U_{i+1,j}^n + B_2 U_{i-1,j}^n + B_3 U_{i,j+1}^n + B_4 U_{i,j-1}^n - \text{Rhs} \right] + (1-\omega) U_{i,j}^n$$

$$\{U\}^{n+1} = \underbrace{-[D]^{-1}[R + S]}_{[G_J]} \{U\}^n + [D]^{-1}\{V\}$$

Jacobi:

$$\{U\}^{n+1} = \underbrace{-[R + D]^{-1}[S]}_{[G_{GS}]} \{U\}^n + [R + D]^{-1}\{V\}$$

Gauss - Seidel:

S.O.R.:

$$\{U\}^{n+1} = \underbrace{[D + \omega R]^{-1} [(1 - \omega)D - \omega S]}_{[G_\omega]} \{U\}^n + [D + \omega R]^{-1} \omega \{V\}$$

Basic Rule:

Type I Boundary: Do not use the PDE on the Bdy.

Type II or III Bdys: Use PDE Plus B.C. together

Spectral Radius, ρ , of iteration matrix $[G]$ is the largest magnitude eigenvalue of $[G]$

$$\rho < 1 \text{ for convergence}$$

Bare Essentials of Iterative Methods

Computational Estimate for ρ

$$\{\delta\}^n = \{U\}^n - \{U\}^{n-1}$$

$$\rho \cong \frac{\|\delta^n\|}{\|\delta^{n-1}\|} = \frac{\left[\sum_{i=1}^M (U_i^n - U_i^{n-1})^2 \right]^{1/2}}{\left[\sum_{i=1}^M (U_i^{n-1} - U_i^{n-2})^2 \right]^{1/2}}$$

Now to prove that an iteration scheme can converge

Consider the following worst case situations.
Recall def. Strict Diagonal Dominance

$$|a_{ii}| > \sum_{j \neq i} |a_{ij}|$$

Expanding on the iteration handout

Define $\{\epsilon\}^n = \{U\}^n - \{U\} = \{U\}^n - [A]^{-1}\{v\}$
 \uparrow exact algebraic solutions (unknown)

Since

$$\{U\}^n = [G]\{U\}^{n-1} + \{r\} \quad \text{and} \quad \{U\} = [G]\{U\} + \{r\}$$

$$\{\epsilon\}^n = [G]\{U\}^{n-1} + \{r\} - G\{U\} - \{r\} = [G](\{U\}^{n-1} - \{U\})$$

or

$$\{\epsilon\}^n = [G]\{\epsilon\}^{n-1} = [G]([G]\{\epsilon\}^{n-2}) \Rightarrow [G]^n\{\epsilon\}^0$$

However we still don't know $\{\epsilon\}^0$

But

Define $\{\delta\}^n = \{U\}^n - \{U\}^{n-1}$ This incremental error can
be determined for all n

$$\{\delta\}^n = [G]\{U\}^{n-1} + \{r\} - [G]\{U\}^{n-2} - \{r\} = [G](\{U\}^{n-1} - \{U\}^{n-2})$$

or

$$\{\delta\}^n = [G]\{\delta\}^{n-1} \Rightarrow [G]^n\{\delta\}^0$$

Finally we can examine Residuals:

$$\begin{aligned} \text{Normally } [A]\{U\} &= \{v\} \\ 0 &= [A]\{U\} - \{v\} \neq [A]\{U\}^n - \{v\} \end{aligned}$$

Define

$$\begin{aligned} \{R\}^n &= [A]\{U\}^n - \{v\} \\ &= [A]\{U\}^n - [A][A]^{-1}\{v\} \\ &= [A](\{U\}^n - [A]^{-1}\{v\}) \\ &\quad \text{remember } \{ \varepsilon \}^n = \{U\}^n - \{U\} \\ \{R\}^n &= [A]\{ \varepsilon \}^n \\ &= [A][G]\{ \varepsilon \}^{n-1} = [A][G][A]^{-1}[A]\{ \varepsilon \}^{n-1} \\ &\quad \{R\}^{n-1} \\ \{R\}^n &= [A][G][A]^{-1}\{R\}^{n-1} \Rightarrow [A][G]^n[A]^{-1}\{R\}^0 \end{aligned}$$

Therefore, we have the following error measures

$$\{ \varepsilon \}^n = [G]\{ \varepsilon \}^{n-1} \Rightarrow [G]^n\{ \varepsilon \}^0 \quad \text{numerical vs algebraic}$$

$$\{ \delta \}^n = [G]\{ \delta \}^{n-1} \Rightarrow [G]^n\{ \delta \}^0 \quad \text{incremental errors}$$

$$\{R\}^n = [G]\{R\}^{n-1} \Rightarrow [G]^n\{R\}^0 \quad \text{residual error}$$

Each of these error indicators converges to zero if and only if the spectral radius, ρ , (or largest absolute value eigenvalue) of the iteration matrix is less than 1.

Therefore,

$$\{ \varepsilon \}^n \cong \rho\{ \varepsilon \}^{n-1}$$

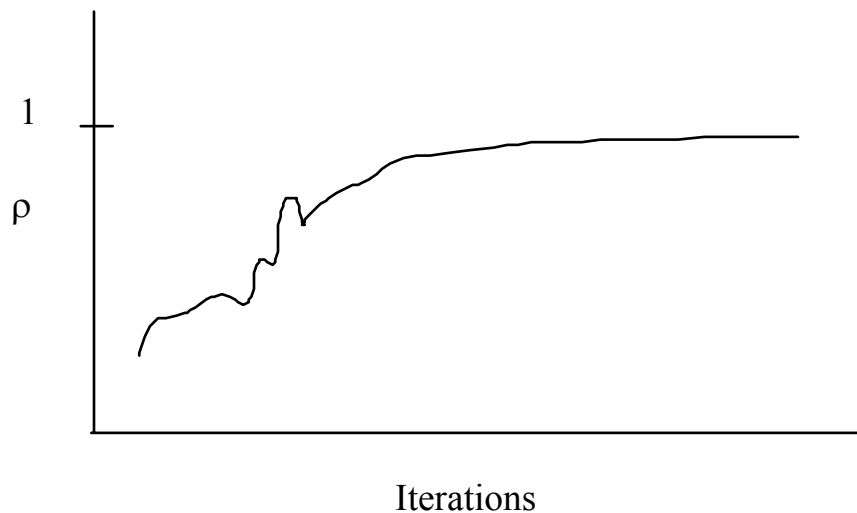
$$\{\delta\}^n \cong \rho \{\delta\}^{n-1}$$

$$\{\mathbf{R}\}^n \cong \rho \{\mathbf{R}\}^{n-1}$$

and one can estimate the spectral radius of the system via

$$\rho \cong \frac{\|\delta^n\|}{\|\delta^{n-1}\|} = \frac{\left[\sum_{i=1}^M (U_i^n - U_i^{n-1})^2 \right]^{1/2}}{\left[\sum_{i=1}^M (U_i^{n-1} - U_i^{n-2})^2 \right]^{1/2}}$$

If one measures ρ expect the following



Given

$$[A] \{U\} = \{v\} \quad \text{where } [A] \text{ strict diagonal dominance}$$

Prove Convergence

Define
$$\theta_i = \sum_{j \neq i} \frac{|a_{ij}|}{|a_{ii}|}$$

Recall
$$\{\varepsilon\}^{n+1} = [G] \{\varepsilon\}^n$$

Jacobi
$$[G_J] = -[D]^{-1}[R + S] = -\frac{\sum_{j \neq i} a_{ij}}{a_{ii}}$$

$$\varepsilon : \varepsilon_i^{n+1} = \frac{1}{a_{ii}} \left[-\sum_{j \neq i} a_{ij} \varepsilon_j^n \right]$$

$$\left| \varepsilon_i^{n+1} \right| \leq \sum_{j \neq i} \frac{|a_{ij}|}{|a_{ii}|} \left| \varepsilon_j^n \right| \leq \theta_i \|\varepsilon\|^n$$

$$\text{where } \|\varepsilon\|^n = \max_j |\varepsilon_j|^n$$

Worst case

$$\|\varepsilon\|^{n+1} \leq \theta_{\max} \|\varepsilon\|^n$$

$$\Rightarrow \theta_{\max} < 1 \quad \text{sufficient for convergence}$$

Note:

Elliptic equation $\theta_{\max} = 1 \quad \therefore$ Jacobi will not diverge.

Examine Gauss - Seidel :

$$\varepsilon_1^{n+1} = \frac{-1}{a_{11}} \sum_{j=2}^N a_{1j} \varepsilon_j^n \Rightarrow \left| \varepsilon_1 \right|^{n+1} \leq \theta_1 \left\| \varepsilon \right\|^n \leq \left\| \varepsilon \right\|^n$$

since $\theta_1 < 1$

$$\varepsilon_2^{n+1} = \frac{-1}{a_{22}} \left[a_{21} \varepsilon_1^{n+1} + \sum_{j=3}^N a_{2j} \varepsilon_j^n \right]$$

$$\left| \varepsilon_2 \right|^{n+1} \leq \frac{|a_{21}|}{|a_{22}|} \left\| \varepsilon \right\|^n + \sum_{j=3}^N \frac{|a_{2j}|}{|a_{22}|} \left| \varepsilon_j \right|^n \leq \theta_2 \left\| \varepsilon \right\|^n$$

etc. . . .

In general since $\frac{|a_{21}|}{|a_{22}|} \left\| \varepsilon_1 \right\|^{n+1} < \frac{|a_{21}|}{|a_{22}|} \left\| \varepsilon \right\|^n$

the Gauss - Seidel will converge faster (or diverge faster) than Jacobi.

and

$$\rho_{GS} = \rho_J^2$$

From S.O.R. theory for [A] Symmetric, Consistently ordered,
 “Property A”

$$\rho_{GS} = \rho_J^2$$

$$\omega_{opt} = \frac{2}{1 + \sqrt{1 - \rho_J^2}} = \frac{2}{1 + \sqrt{1 - \rho_{GS}}}$$

Recall: Self Adjoint
implies symmetry

Rates of Convergence

In the limit of large n, recall that

$$\{\epsilon\}^{n+M} = \rho^M \{\epsilon\}^n \quad \text{and therefore}$$

$$\|\epsilon\|^{n+M} = \rho^M \|\epsilon\|^n$$

or

$$\rho^M = \frac{\|\epsilon\|^{n+M}}{\|\epsilon\|^n}$$

and if you wish to reduce the existing error by a factor of K

$$K = \frac{\|\epsilon\|^{n+M}}{\|\epsilon\|^n}$$

one can write

$\rho^M = K$ and solve for M, the number of iterations
 required to get to desired accuracy

$$M = \ln(K) / \ln(\rho)$$

When solving $\nabla^2 U = 0$ on a square

Jacobi

$$\rho = \max|\lambda| = \cos h \sim 1 - \frac{h^2}{2} \text{ as } h \rightarrow 0$$

The rate of convergence of a linear iteration

$$\{U\}^{n+1} = [G]\{U\}^n + \{r\} \quad \text{characterized by the matrix } [G] \text{ is}$$

$$R(G_J) = -\log \lambda(G_J) = -\log \rho$$

(-) since $\rho < 1$ for (+) convergence rate

Ref:

Young, D.M. Trans. AM. Math. Soc., 76, #92, 1954

Ames - on reserve

Westlake - listed in handout class 1 - Appendix B

Eigenvalue Bounds

and $-\log \rho \sim -\log\left(1 - \frac{h^2}{2}\right) = \frac{h^2}{2} + O(h^4)$

Thus, the convergence rate for Jacobi iterations is approximately $h^2/2$ which is slow for small values of h

$$\log(1 + x) = x - \frac{x^2}{2} + \frac{x^3}{3} - \frac{x^4}{4} \dots \quad -1 < x \leq 1$$

The $\max |\lambda|$ in $[G_{GS}]$ is $\cos^2 h \sim 1 - h^2$ as $h \rightarrow 0$

$$\therefore R(G_{GS}) = -\log(\cos^2 h) \sim h^2 + O(h^4)$$

or the Gauss-Seidel iterations will converge twice as fast as Jacobi.

Finally

$$\max |\lambda| \text{ in } [G_{\omega}] \sim 1 - 2h \text{ as } h \rightarrow 0$$

$$R(G) \sim 2h + O(h^2) \text{ for optimal S.O.R.}$$

$$\frac{2h}{h^2} \sim \frac{2}{h} \text{ times faster than Gauss-Seidel which for}$$

small h is significant